Random networks

Pure, abstract random networks:

- Consider set of all networks with \( N \) labelled nodes and \( m \) edges.
- Standard random network = randomly chosen network from this set.
- To be clear: each network is equally probable.
- Sometimes equiprobability is a good assumption, but it is always an assumption.
- Known as Erdős-Rényi random networks or ER graphs.
Random networks

How to build standard random networks:

- Given $N$ and $m$.
- Two probabilistic methods (we'll see a third later on)

1. Connect each of the $\binom{N}{2}$ pairs with appropriate probability $p$.
   - Useful for theoretical work.
2. Take $N$ nodes and add exactly $m$ links by selecting edges without replacement.
   - Algorithm: Randomly choose a pair of nodes $i$ and $j$, $i \neq j$, and connect if unconnected; repeat until all $m$ edges are allocated.
   - Best for adding small numbers of links (most cases).
   - 1 and 2 are effectively equivalent for large $N$.

A few more things:

- For method 1, # links is probabilistic:
  $$\langle m \rangle = p \binom{N}{2} = p \frac{1}{2} N(N - 1)$$
- So the expected or average degree is
  $$\langle k \rangle = \frac{2 \langle m \rangle}{N} = \frac{2}{N} p \frac{1}{2} N(N - 1) = p(N - 1).$$
- Which is what it should be...
- If we keep $\langle k \rangle$ constant then $p \propto 1/N \to 0$ as $N \to \infty$.

Random networks: examples for $N=500$

Random networks: largest components
Random networks: examples for $N=500$

For method 1, what is the clustering coefficient for a finite network?

Consider triangle/triple clustering coefficient (Newman [1]):

$$C_2 = \frac{3 \times \text{#triangles}}{\text{#triples}}$$

Recall: $C_2 =$ probability that two nodes are connected given they have a friend in common.

For standard random networks, we have simply that $C_2 = p$.

Random networks: largest components

So for large random networks ($N \to \infty$), clustering drops to zero.

Key structural feature of random networks is that they locally look like branching networks (no loops).
Random networks

Degree distribution:

- Recall $p_k = \text{probability that a randomly selected node has degree } k$.
- Consider method 1 for constructing random networks: each possible link is realized with probability $p$.
- Now consider one node: there are $\binom{N}{k}$ ways the node can be connected to $k$ of the other $N - 1$ nodes.
- Each connection occurs with probability $p$, each non-connection with probability $(1 - p)$.
- Therefore have a binomial distribution:

$$P(k; p, N) = \binom{N - 1}{k} p^k (1 - p)^{N - 1 - k}.$$  

Random networks

Limiting form of $P(k; p, N)$:

- Substitute $p = \frac{\langle k \rangle}{N - 1}$ into $P(k; p, N)$ and hold $k$ fixed:

$$P(k; p, N) = \binom{N - 1}{k} \left( \frac{\langle k \rangle}{N - 1} \right)^k \left( 1 - \frac{\langle k \rangle}{N - 1} \right)^{N - 1 - k}$$

$$= \frac{(N - 1)!}{k!(N - 1 - k)!} \left( \frac{\langle k \rangle}{N - 1} \right)^k \left( 1 - \frac{\langle k \rangle}{N - 1} \right)^{N - 1 - k}$$

$$= \frac{(N - 1)(N - 2) \cdots (N - k)}{k!} \left( \frac{\langle k \rangle}{N - 1} \right)^k \left( 1 - \frac{\langle k \rangle}{N - 1} \right)^{N - 1 - k}$$

- Now use the excellent result:

$$P(k; p, N) \approx \frac{\langle k \rangle^k}{k!} \left( 1 - \frac{\langle k \rangle}{N - 1} \right)^{N - 1 - k}$$

- Identifying $n = N - 1$ and $x = -\langle k \rangle$:

$$P(k; \langle k \rangle) \approx \frac{\langle k \rangle^k}{k!} e^{-\langle k \rangle} \left( 1 - \frac{\langle k \rangle}{N - 1} \right)^{-k} \frac{\langle k \rangle^k}{k!} e^{-\langle k \rangle}$$

- This is a Poisson distribution with mean $\langle k \rangle$.  

References

Generating Functions

Limiting form of $P(k; p, N)$:

- We are now here:

$$P(k; p, N) \approx \frac{\langle k \rangle^k}{k!} \left( 1 - \frac{\langle k \rangle}{N - 1} \right)^{N - 1 - k}$$

- Now use the excellent result:

$$\lim_{n \to \infty} \left( 1 + \frac{x}{n} \right)^n = e^x.$$  

(Use l'Hôpital's rule to prove.)

- Identifying $n = N - 1$ and $x = -\langle k \rangle$:

$$P(k; \langle k \rangle) \approx \frac{\langle k \rangle^k}{k!} e^{-\langle k \rangle} \left( 1 - \frac{\langle k \rangle}{N - 1} \right)^{-k} \frac{\langle k \rangle^k}{k!} e^{-\langle k \rangle}$$

- This is a Poisson distribution with mean $\langle k \rangle$.  

References

Generating Functions
General random networks

- So... standard random networks have a Poisson degree distribution
- Generalize to arbitrary degree distribution $P_k$.
- Also known as the configuration model[1].
- Can generalize construction method from ER random networks.
- Assign each node a weight $w$ from some distribution $P_w$ and form links with probability $P(\text{link between } i \text{ and } j) \propto w_i w_j$.
- But we’ll be more interested in
  1. Randomly wiring up (and rewiring) already existing nodes with fixed degrees.
  2. Examining mechanisms that lead to networks with certain degree distributions.

Random networks: largest components

Poisson basics:

- Normalization: we must have
  $$\sum_{k=0}^{\infty} P(k; \langle k \rangle) = 1$$
- Checking:
  $$\sum_{k=0}^{\infty} P(k; \langle k \rangle) = \sum_{k=0}^{\infty} \frac{\langle k \rangle^k}{k!} e^{-\langle k \rangle}$$
  $$= e^{-\langle k \rangle} \sum_{k=0}^{\infty} \frac{\langle k \rangle^k}{k!}$$
  $$= e^{-\langle k \rangle} e^{\langle k \rangle} = 1 \checkmark$$
Poisson basics:

- Mean degree: we must have
  \[ \langle k \rangle = \sum_{k=0}^{\infty} k P(k; \langle k \rangle). \]

- Checking:
  \[ \sum_{k=0}^{\infty} k P(k; \langle k \rangle) = \sum_{k=0}^{\infty} k \frac{k^k}{k!} e^{-k} \]
  \[ = e^{-k} \sum_{k=0}^{\infty} \frac{k^k}{(k-1)!} \]
  \[ = \langle k \rangle e^{-k} \sum_{k=0}^{\infty} \frac{k^{k-1}}{(k-1)!} \]
  \[ = \langle k \rangle e^{-k} (\sum_{i=0}^{\infty} \frac{k^i}{i!}) = \langle k \rangle e^{-k} e^{\langle k \rangle} = \langle k \rangle \]

- We'll get to a better way of doing this...

The edge-degree distribution:

- The degree distribution \( P_k \) is fundamental for our description of many complex networks
- Again: \( P_k \) is the degree of a randomly chosen node.
- A second very important distribution arises from choosing randomly on edges rather than on nodes.
- Define \( Q_k \) to be the probability the node at a random end of a randomly chosen edge has degree \( k \).
- Now choosing nodes based on their degree (i.e., size):
  \[ Q_k \propto k P_k \]

- Normalized form:
  \[ Q_k = \frac{k P_k}{\sum_{k'=0}^{\infty} k' P_{k'}} = \frac{k P_k}{\langle k \rangle}. \]

The edge-degree distribution:

- For random networks, \( Q_k \) is also the probability that a friend (neighbor) of a random node has \( k \) friends.
- Useful variant on \( Q_k \):
  \[ R_k = \text{probability that a friend of a random node has \( k \) other friends}. \]
  \[ R_k = \frac{(k+1) P_{k+1}}{\sum_{k'=0}^{\infty} (k' + 1) P_{k' + 1}} = \frac{(k+1) P_{k+1}}{\langle k \rangle}. \]

- Equivalent to friend having degree \( k + 1 \).
- **Natural question:** what’s the expected number of other friends that one friend has?
The edge-degree distribution:

- Given $R_k$ is the probability that a friend has $k$ other friends, then the average number of friends' other friends is

$$\langle k \rangle_R = \sum_{k=0}^{\infty} k R_k = \sum_{k=0}^{\infty} k \frac{(k+1) P_{k+1}}{\langle k \rangle}$$

$$= \frac{1}{\langle k \rangle} \sum_{k=1}^{\infty} k(k+1) P_{k+1}$$

$$= \frac{1}{\langle k \rangle} \sum_{k=1}^{\infty} \left((k+1)^2 - (k+1)\right) P_{k+1}$$

(where we have sneakily matched up indices)

$$= \frac{1}{\langle k \rangle} \sum_{j=0}^{\infty} (j^2 - j) P_j \quad \text{(using } j = k+1)$$

$$= \frac{1}{\langle k \rangle} \left(\langle k^2 \rangle - \langle k \rangle\right)$$

Two reasons why this matters

**Reason #1:**
- Average # friends of friends per node is

$$\langle k_2 \rangle = \langle k \rangle \times \langle k \rangle_R = \langle k \rangle \left(\frac{1}{\langle k \rangle} \left(\langle k^2 \rangle - \langle k \rangle\right)\right) = \langle k^2 \rangle - \langle k \rangle.$$

- Key: Average depends on the 1st and 2nd moments of $P_k$ and not just the 1st moment.
- Three peculiarities:
  1. We might guess $\langle k_2 \rangle = \langle k \rangle(\langle k \rangle - 1)$ but it’s actually $\langle k(k-1) \rangle$.
  2. If $P_k$ has a large second moment, then $\langle k_2 \rangle$ will be big.
     (e.g., in the case of a power-law distribution)
  3. Your friends are different to you...

Two reasons why this matters

**Reason #1:**
- A node’s average # of friends: $\langle k \rangle$
- Friend’s average # of friends: $\langle k^2 \rangle$  
- Comparison:

$$\frac{\langle k^2 \rangle}{\langle k \rangle^2} = \langle k \rangle \frac{\langle k^2 \rangle}{\langle k \rangle^2} = \langle k \rangle \frac{\sigma^2 + \langle k \rangle^2}{\langle k \rangle^2} = \langle k \rangle \left(1 + \frac{\sigma^2}{\langle k \rangle^2}\right) \geq \langle k \rangle$$

- So only if everyone has the same degree (variance $= \sigma^2 = 0$) can a node be the same as its friends.
- Intuition: for random networks, the more connected a node, the more likely it is to be chosen as a friend.
Two reasons why this matters

(Big) Reason #2:

- $\langle k \rangle_R$ is key to understanding how well random networks are connected together.
- e.g., we’d like to know what’s the size of the largest component within a network.
- As $N \to \infty$, does our network have a giant component?
- Defn: Component = connected subnetwork of nodes such that $\exists$ path between each pair of nodes in the subnetwork, and no node out side of the subnetwork is connected to it.
- Defn: Giant component = component that comprises a non-zero fraction of a network as $N \to \infty$.
- Note: Component = Cluster

Giant component

Standard random networks:

- Recall $\langle k^2 \rangle = \langle k \rangle^2 + \langle k \rangle$.
- Condition for giant component:
  $$\langle k \rangle_R = \frac{\langle k^2 \rangle - \langle k \rangle}{\langle k \rangle} = \frac{\langle k^2 \rangle}{\langle k \rangle} - \frac{\langle k \rangle}{\langle k \rangle} = \langle k \rangle$$
  
  Therefore when $\langle k \rangle > 1$, standard random networks have a giant component.

- When $\langle k \rangle < 1$, all components are finite.
- Fine example of a continuous phase transition (H).  
- We say $\langle k \rangle = 1$ marks the critical point of the system.

Random networks with skewed $P_k$:

- e.g., if $P_k = c k^{-\gamma}$ with $2 < \gamma < 3$ then
  
  $$\langle k^2 \rangle = c \sum_{k=0}^{\infty} k^2 k^{-\gamma} \sim \int_{x=0}^{\infty} x^{2-\gamma} dx \propto x^{3-\gamma}\bigg|_{x=0}^{\infty} = \infty \quad (> \langle k \rangle).$$

  So giant component always exists for these kinds of networks.

  Cutoff scaling is $k^{-3}$: if $\gamma > 3$ then we have to look harder at $\langle k \rangle_R$.  

Structure of random networks

Giant component:

- A giant component exists if when we follow a random edge, we are likely to hit a node with at least 1 other outgoing edge.

- Equivalently, expect exponential growth in node number as we move out from a random node.

- All of this is the same as requiring $\langle k \rangle_R > 1$.

- Giant component condition (or percolation condition):
  $$\langle k \rangle_R = \frac{\langle k^2 \rangle - \langle k \rangle}{\langle k \rangle} > 1$$

  Again, see that the second moment is an essential part of the story.

  Equivalent statement: $\langle k^2 \rangle > 2 \langle k \rangle$
Giant component

And how big is the largest component?

- Define $S_1$ as the size of the largest component.
- Consider an infinite ER random network with average degree $\langle k \rangle$.
- Let’s find $S_1$ with a back-of-the-envelope argument.
- Define $\delta$ as the probability that a randomly chosen node does not belong to the largest component.
- Simple connection: $\delta = 1 - S_1$.
- Dirty trick: If a randomly chosen node is not part of the largest component, then none of its neighbors are.
- So
  $$\delta = \sum_{k=0}^{\infty} P_k \delta^k$$
- Substitute in Poisson distribution...

We can figure out some limits and details for $S_1 = 1 - e^{-\langle k \rangle S_1}$.

- As $\langle k \rangle \to 0$, $S_1 \to 0$.
- As $\langle k \rangle \to \infty$, $S_1 \to 1$.
- Notice that at $\langle k \rangle = 1$, the critical point, $S_1 = 0$.
- Only solvable for $S > 0$ when $\langle k \rangle > 1$.
- Really a transcritical bifurcation \[^2\].

Now substitute in $\delta = 1 - S_1$ and rearrange to obtain a transcendental equation for $S_1$:

$$S_1 = 1 - e^{-\langle k \rangle S_1}.$$
Giant component

Turns out we were lucky...

- Our dirty trick only works for ER random networks.
- The problem: We assumed that neighbors have the same probability $\delta$ of belonging to the largest component.
- But we know our friends are different from us...
- Works for ER random networks because $\langle k \rangle = \langle k \rangle_R$.
- We need a separate probability $\delta'$ for the chance that a node at the end of a random edge is part of the largest component.
- We can do this but we need to enhance our toolkit with Generatingfunctionology... \[3\]

Example

- Take a degree distribution with exponential decay:
  
  \[ P_k = c e^{-\lambda k} \]

  where $c = 1 - e^{-\lambda}$.

  - The generating function for this distribution is
    
    \[ F(x) = \sum_{k=0}^{\infty} P_k x^k = \sum_{k=0}^{\infty} c e^{-\lambda k} x^k = \frac{c}{1 - xe^{-\lambda}}. \]

- Notice that $F(1) = c/(1 - e^{-\lambda}) = 1$.

- For probability distributions, we must always have $F(1) = 1$ since

  \[ F(1) = \sum_{k=0}^{\infty} P_k 1^k = \sum_{k=0}^{\infty} P_k = 1. \]

Generating functions

- Idea: Given a sequence $a_0, a_1, a_2, \ldots$, associate each element with a distinct function or other mathematical object.
- Well-chosen functions allow us to manipulate sequences and retrieve sequence elements.

Definition:

- The generating function (g.f.) for a sequence $\{a_n\}$ is

  \[ F(x) = \sum_{n=0}^{\infty} a_n x^n. \]

- Roughly: transforms a vector in $R^\infty$ into a function defined on $R^1$.
- Related to Fourier, Laplace, Mellin, . . .

Properties of generating functions

- Average degree:

  \[ \langle k \rangle = \sum_{k=0}^{\infty} kP_k = \sum_{k=0}^{\infty} kP_k x^{k-1} \bigg|_{x=1} = \frac{d}{dx} F(x) \bigg|_{x=1} = F'(1) \]

- In general, many calculations become simple, if a little abstract.

- For our exponential example:

  \[ F'(x) = \frac{(1 - e^{-\lambda}) e^{-\lambda}}{(1 - xe^{-\lambda})^2}. \]

- So:

  \[ \langle k \rangle = F'(1) = \frac{e^{-\lambda}}{(1 - e^{-\lambda})}. \]
Properties of generating functions

Useful pieces for probability distributions:

- Normalization:
  \[ F(1) = 1 \]

- First moment:
  \[ \langle k \rangle = F'(1) \]

- Higher moments:
  \[ \langle k^n \rangle = \left. x^n \frac{d^n}{dx^n} F(x) \right|_{x=1} \]

- \( k \)th element of sequence (general):
  \[ P_k = \left. \frac{1}{k!} \frac{d^k}{dx^k} F(x) \right|_{x=0} \]

Edge-degree distribution

We have

\[
F_R(x) = \sum_{k=0}^{\infty} R_k x^k = \sum_{k=0}^{\infty} \frac{(k+1)P_{k+1}}{\langle k \rangle} x^k.
\]

Shift index to \( j = k + 1 \) and pull out \( \frac{1}{\langle k \rangle} \):

\[
F_R(x) = \frac{1}{\langle k \rangle} \sum_{j=1}^{\infty} jP_j x^{j-1} = \frac{1}{\langle k \rangle} \sum_{j=0}^{\infty} P_j j x^j
\]

\[
= \frac{1}{\langle k \rangle} \sum_{j=0}^{\infty} \frac{d}{dx} P_j x^j = \frac{1}{\langle k \rangle} \frac{d}{dx} \sum_{j=0}^{\infty} P_j x^j = \frac{1}{\langle k \rangle} F'_P(x).
\]

Finally, since \( \langle k \rangle = F'_P(1) \),

\[
F_R(x) = \frac{F'_P(x)}{F'_P(1)}.
\]

Edge-degree distribution

Recall our condition for a giant component:

\[ \langle k \rangle_R = \frac{\langle k^2 \rangle - \langle k \rangle}{\langle k \rangle} > 1. \]

Let’s reexpress our condition in terms of generating functions.

- We first need the g.f. for \( R_k \).
- We’ll now use this notation:
  \( F_P(x) \) is the g.f. for \( P_k \).
  \( F_R(x) \) is the g.f. for \( R_k \).
- Condition in terms of g.f. is:
  \[ \langle k \rangle_R = F'_R(1) > 1. \]
- Now find how \( F_R \) is related to \( F_P \ldots \)

Recall giant component condition is

\[ \langle k \rangle_R = F'_R(1) > 1. \]

Since we have \( F_R(x) = F'_P(x) F'_R(1) \),

\[ F'_R(x) = \frac{F''_P(x)}{F'_P(1)}. \]

Setting \( x = 1 \), our condition becomes

\[ \frac{F''_P(1)}{F'_P(1)} > 1. \]
Size distributions

To figure out the size of the largest component \((S_1)\), we need more resolution on component sizes.

Definitions:

- \(\pi_n\) = probability that a random node belongs to a finite component of size \(n < \infty\).
- \(\rho_n\) = probability that a random link leads to a finite subcomponent of size \(n < \infty\).

Local-global connection:

\[
P_k, R_k \leftrightarrow \pi_n, \rho_n
\]

neighbors \(\leftrightarrow\) components

Useful results we'll need for g.f.'s

Sneaky Result 1:

- Consider two random variables \(U\) and \(V\) whose values may be 0, 1, 2, \ldots
- Write probability distributions as \(U_k\) and \(V_k\) and g.f.'s as \(F_U\) and \(F_V\).
- \(SR1\): If a third random variable is defined as

\[
W = \sum_{i=1}^{V} U^{(i)} \; \text{with each} \; U^{(i)} \overset{d}{=} U
\]

then

\[
F_W(x) = F_V(F_U(x))
\]

Proof of SN1:

Write probability that variable \(W\) has value \(k\) as \(W_k\).

\[
W_k = \sum_{j=0}^{\infty} V_j \times \Pr(\text{sum of } j \text{ draws of variable } U = k)
\]

\[
= \sum_{j=0}^{\infty} V_j \sum_{\{i_1, i_2, \ldots, i_k\}} \prod_{h=1}^{j} U_{i_h} \cdot \cdots \cdot U_j
\]

\[
\Rightarrow F_W(x) = \sum_{k=0}^{\infty} W_k x^k = \sum_{k=0}^{\infty} \sum_{j=0}^{\infty} V_j \sum_{\{i_1, i_2, \ldots, i_k\}} \prod_{h=1}^{j} U_{i_h} \cdot \cdots \cdot U_j x^k
\]

\[
= \sum_{j=0}^{\infty} V_j \sum_{k=0}^{\infty} \sum_{\{i_1, i_2, \ldots, i_k\}} \prod_{h=1}^{j} U_{i_h} x^{i_1} U_{i_2} x^{i_2} \cdots U_j x^j
\]
Proof of SN1:

With some concentration, observe:

\[ F_W(x) = \sum_{j=0}^{\infty} V_j \sum_{k=0}^{\infty} \sum_{(l_1, l_2, \ldots, l_k) \mid l_1 + l_2 + \ldots + l_k = k} U_{l_1} x^{l_1} U_{l_2} x^{l_2} \cdots U_{l_k} x^{l_k} \]

\( x^k \) piece of \( \left( \sum_{j=0}^{\infty} U_j x^j \right)^k \)

\( \left( \sum_{j=0}^{\infty} U_j x^j \right)^k = (F_U(x))^k \)

\[ = \sum_{j=0}^{\infty} V_j (F_U(x))^j \]

\[ = F_V(F_U(x)) \checkmark \]

Useful results we’ll need for g.f.’s

Sneaky Result 2:

- Start with a random variable \( U \) with distribution \( U_k \) \( (k = 0, 1, 2, \ldots) \)
- \( \text{SNR2: If a second random variable is defined as} \)
  \[ V = U + 1 \text{ then } F_V(x) = xF_U(x) \]

- Reason: \( V_k = U_{k-1} \) for \( k \geq 1 \) and \( V_0 = 0 \).

\[ \therefore F_V(x) = \sum_{k=0}^{\infty} V_k x^k = \sum_{k=1}^{\infty} U_{k-1} x^k \]

\[ = x \sum_{j=0}^{\infty} U_j x^j = xF_U(x). \checkmark \]

Connecting generating functions

- Goal: figure out forms of the component generating functions, \( F_\pi \) and \( F_\rho \).
- \( \pi_n \) = probability that a random node belongs to a finite component of size \( n \)
  \[ = \sum_{k=0}^{\infty} P_k \times \Pr \left( \text{sum of sizes of subcomponents at end of } k \text{ random links } \approx n - 1 \right) \]

\[ \therefore \]

\[ F_\pi(x) = \frac{\chi_{SN2}}{F_{SN1}(F_\rho(x))} \]

- Extra factor of \( x \) accounts for random node itself.
Connecting generating functions

- \( \rho_n \) = probability that a random link leads to a finite subcomponent of size \( n \).
- Invoke one step of recursion: \( \rho_n \) = probability that a random node arrived along a random edge is part of a finite subcomponent of size \( n \).

\[
\sum_{k=0}^{\infty} R_k \times \Pr \left( \text{sum of sizes of subcomponents at end of } k \text{ random links } = n - 1 \right)
\]

Therefore:

\[
F_\rho(x) = \frac{1}{\rho} \left( \frac{F_R(F_\rho(x))}{F_N} \right)
\]

Again, extra factor of \( x \) accounts for random node itself.

Component sizes

- Remembering vaguely what we are doing:
  - Finding \( F_P \) to obtain the size of the largest component \( S_1 = 1 - F_\pi(1) \).
  - Set \( x = 1 \) in our two equations:
    \[
    F_\pi(1) = F_P(F_\rho(1)) \quad \text{and} \quad F_\rho(1) = F_R(F_\rho(1))
    \]
  - Solve second equation numerically for \( F_\rho(1) \).
  - Plug \( F_\rho(1) \) into first equation to obtain \( F_\pi(1) \).

Connecting generating functions

- We now have two functional equations connecting our generating functions:
  \[
  F_\pi(x) = xF_P(F_\rho(x)) \quad \text{and} \quad F_\rho(x) = xF_R(F_\rho(x))
  \]
- Taking stock: We know \( F_P(x) \) and \( F_R(x) = F'_P(x)/F'_P(1) \).
- We first untangle the second equation to find \( F_\rho \).
- We can do this because it only involves \( F_\rho \) and \( F_R \).
- The first equation then immediately gives us \( F_\pi \) in terms of \( F_\rho \) and \( F_R \).

Example: Standard random graphs.

- We can show \( F_P(x) = e^{-(k)(1-x)} \)
  \[
  \therefore \quad F_R(x) = F'_P(x)/F'_P(1) = e^{-(k)(1-x')}/e^{-(k)(1-x')}|_{x'=1} = e^{-(k)(1-x)} = F_P(x) \quad \text{...aha!}
  \]
- RHS’s of our two equations are the same.
- So \( F_\pi(x) = F_\rho(x) = xF_R(F_\rho(x)) = xF_R(F_\pi(x)) \)
- Why our dirty (but wrong) trick worked earlier...
Component sizes

- We are down to \( F_\pi(x) = xF_R(F_\pi(x)) \) and \( F_R(x) = xe^{-(k)(1-x)} \).
- \( \therefore F_\pi(x) = xe^{-(k)(1-F_\pi(x))} \)
- We’re first after \( S_1 = 1 - F_\pi(1) \) so set \( x = 1 \) and replace \( F_\pi(1) \) by \( 1 - S_1 \):
  \[
  1 - S_1 = e^{-(k)S_1}
  \]
- Just as we found with our dirty trick...
- Again, have to resort to numerics at this point.

Average component size

**Example:** Standard random graphs.

- Use fact that \( F_P = F_R \) and \( F_\pi = F_\rho \).
- Two differentiated equations reduce to only one:
  \[
  F'_\pi(x) = F_P(F_\pi(x)) + xF'_\pi(x)F_P(F_\pi(x))
  \]
  Rearrange: \( F'_\pi(x) = \frac{F_P(F_\pi(x))}{1 - xF'_P(F_\pi(x))} \)
- Simplify denominator using \( F'_P(x) = \langle k \rangle F_\pi(x) \)
- Replace \( F_P(F_\pi(x)) \) using \( F_\pi(x) = xF_P(F_\pi(x)) \).
- Set \( x = 1 \) and replace \( F_\pi(1) \) with \( 1 - S_1 \).
- End result: \( \langle n \rangle = F'_\pi(1) = \frac{(1 - S_1)}{1 - \langle k \rangle(1 - S_1)} \)

Average component size

- Next: find average size of finite components \( \langle n \rangle \).
- Using standard G.F. result: \( \langle n \rangle = F'_\pi(1) \).
- Try to avoid finding \( F_\pi(x) \)...
- Starting from \( F_\pi(x) = xF_P(F_\pi(x)) \), we differentiate:
  \[
  F'_\pi(x) = F_P(F_\pi(x)) + xF'_P(x)F_P(F_\pi(x))
  \]
- While \( F_P(x) = xF_R(F_\rho(x)) \) gives
  \[
  F'_\rho(x) = F_R(F_\rho(x)) + xF'_\rho(x)F_R(F_\rho(x))
  \]
- Now set \( x = 1 \) in both equations.
- We solve the second equation for \( F'_\rho(1) \) (we must already have \( F_\rho(1) \)).
- Plug \( F'_\rho(1) \) and \( F_\rho(1) \) into first equation to find \( F'_\pi(1) \).

Average component size

- Our result for standard random networks:
  \[
  \langle n \rangle = F'_\pi(1) = \frac{(1 - S_1)}{1 - \langle k \rangle(1 - S_1)}
  \]
- Recall that \( \langle k \rangle = 1 \) is the critical value of average degree for standard random networks.
- Look at what happens when we increase \( \langle k \rangle \) to 1 from below.
- We have \( S_1 = 0 \) for all \( \langle k \rangle < 1 \) so
  \[
  \langle n \rangle = \frac{1}{1 - \langle k \rangle}
  \]
- This blows up as \( \langle k \rangle \rightarrow 1 \).
- **Reason:** we have a power law distribution of component sizes at \( \langle k \rangle = 1 \).
- Typical critical point behavior....
Average component size

- Limits of \( \langle k \rangle = 0 \) and \( \infty \) make sense for
  \[
  \langle n \rangle = F'_{\pi}(1) = \frac{(1 - S_1)}{1 - \langle k \rangle (1 - S_1)}
  \]
- As \( \langle k \rangle \to 0 \), \( S_1 = 0 \), and \( \langle n \rangle \to 1 \).
- All nodes are isolated.
- As \( \langle k \rangle \to \infty \), \( S_1 \to 1 \) and \( \langle n \rangle \to 0 \).
- No nodes are outside of the giant component.

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